CS631 - Advanced Programming in the UNIX Environment

File Systems, System Data Files, Time & Date

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Lecture 04: File Systems, System Data Files, Time & Date

September 23, 2019

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- you can have many such mappings to the same file



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- each directory contains at least two entries:
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- the link count (st_nlink) of a directory is at least 2



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		Returns: 0 if OK, -1 on error

- Creates a link to an existing file (hard link).
- POSIX.1 allows links to cross filesystems, most implementations (SVR4, BSD) don't.
- only uid(0) can create links to directories (loops in filesystem are bad)

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<pre>#include <unistd.h></unistd.h></pre>	
<pre>int unlink(const char *path);</pre>	
	Returns: 0 if OK, -1 on error

- removes directory entry and decrements link count of file
- if file link count == 0, free data blocks associated with file (...unless processes have the file open)

link(2) and unlink(2)

- \$ cc -Wall wait-unlink.c
- \$./a.out
- \$ df .



#include <stdio.h>
int rename(const char *from, const char *to);
Returns: 0 if OK, -1 on error

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Symbolic Links

#include <unistd.h>

int symlink(const char *name1, const char *name2);

Returns: 0 if OK, -1 on error

- file whose "data" is a path to another file
- anyone can create symlinks to directories or files
- certain functions dereference the link, others operate on the link

How do we get the contents of a symlink? open(2) and read(2)?

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#include <unistd.h>

int readlink(const char *path, char *buf, size_t bufsize);

Returns: number of bytes placed into buffer if OK, -1 on error

This function combines the actions of open, read, and close. Note: *buf* is not NUL terminated.

File Times

#include <sys/types.h>
int utimes(const char *path, const struct timeval times[2]);
int lutimes(const char *path, const struct timeval times[2]);
int futimes(int fd, const struct timeval times[2]);
Returns: 0 if OK, -1 on error

If *times* is NULL, access time and modification time are set to the current time (must be owner of file or have write permission). If *times* is non-NULL, then times are set according to the timeval struct array. For this, you must be the owner of the file (write permission not enough).

Note that st_ctime is set to the current time in both cases.

For the effect of various functions on the access, modification and changes-status times see Stevens, p. 117.

Note: some systems implement lutimes(3) (library call) via utimes(2) syscalls.

mkdir(2) and rmdir(2)

#include <sys/types.h>
#include <sys/stat.h>

int mkdir(const char *path, mode_t mode);

Returns: 0 if OK, -1 on error

Creates a new, empty (except for . and .. entries) directory. Access permissions specified by *mode* and restricted by the umask(2) of the calling process.

#include <unistd.h>
int rmdir(const char *path);
Returns: 0 if OK, -1 on error

If the link count is 0 (after this call), and no other process has the directory open, directory is removed. Directory must be empty (only . and .. remaining)

Reading Directories

- read by anyone with read permission on the directory
- format of directory is implementation dependent (always use readdir and friends)

opendir, readdir and closedir should be familiar from our small 1s clone. rewinddir resets an open directory to the beginning so readdir will again return the first entry.

For directory traversal, consider fts(3) (not available on all UNIX versions).

Moving around directories

#include <unistd.h>

char *getcwd(char *buf, size_t size);

Returns: buf if OK, NULL on error

Returns: 0 if OK, -1 on error

Get the kernel's idea of our process's current working directory.

<pre>#include <unistd.h></unistd.h></pre>	
<pre>int chdir(const char *path); int fchdir(int fd);</pre>	

Allows a process to change its current working directory. Note that chdir and fchdir affect only the current process.

\$ cc -Wall cd.c
\$./a.out /tmp

Password File

Called a *user database* by POSIX and usually found in /etc/passwd, the password file contains the following fields:

Description	struct passwd member	POSIX.1
username	char *pw_name	Х
hashed passwd	char *pw_passwd	
numerical user id	uid_t pw_uid	Х
numerical group id	gid_t pw_gid	Х
comment field	char *pw_gecos	
initial working directory	char *pw_dir	Х
initial shell	char *pw_shell	Х

Password field is a one-way hash of the users password. Some fields can be empty:

- password empty implies no password
- shell empty implies /bin/sh

Password File

```
#include <sys/types.h>
#include <pwd.h>
struct passwd *getpwuid(uid_t uid);
```

struct passwd *getpwnam(const char *name);

Returns: pointer if OK, NULL on error

- getpwent returns next password entry in file each time it's called, no order
- setpwent rewinds to "beginning" of entries
- endpwent closes the file(s)

```
See also: getspnam(3)/getspent(3) (where available)
```

Group File

Called a *group database* by POSIX and usually found in /etc/group, the group file contains the following fields:

Description	struct group member	POSIX.1
groupname	char *gr_name	Х
hashed passwd	char *gr_passwd	
numerical group id	uid_t gr_uid	Х
array of pointers to user names	char **gr_mem	Х

The gr_mem array is terminated by a NULL pointer.

Group File

#include <sys/types.h>
#include <grp.h>

struct group *getgrgid(gid_t gid);
struct group *getgrnam(const char *name);

Returns: pointer if OK, NULL on error

These allow us to look up an entry given a user's group name or numerical GID. What if we need to go through the group file entry by entry? Nothing in POSIX.1, but SVR4 and BSD give us:

<pre>#include <sys types.h=""> #include <grp.h></grp.h></sys></pre>	
<pre>struct group *getgrent(void);</pre>	Returns: pointer if OK, NULL on error
<pre>void setgrent(void); void endgrent(void);</pre>	

- getgrent returns next group entry in file each time it's called, no order
- setgrent rewinds to "beginning" of entries
- endgrent closes the file(s)

Supplementary Groups and other data files

#include <sys/types.h>
#include <unistd.h>

Note: if gidsetsize == 0, getgroups(2) returns number of groups without modifying grouplist.

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Other system databases

Similar routines as for password/group for accessing system data files:

Description	Data file	Header	Structure	Additional lookup functions
hosts	/etc/hosts	<netbdb.h></netbdb.h>	hostent	gethostbyname
				gethostbyaddr
networks	/etc/networks	<netbdb.h>	netent	genetbyname
				getnetbyaddr
protocols	/etc/protocols	<netbdb.h>	protoent	getprotobyname
				getprotobynumber
services	/etc/services	<netbdb.h>	servent	getservbyname
				getservbyport

System Identification

- Pass a pointer to a utsname struct. This struct contains fields like opsys name, version, release, architecture, etc.
- This function used by the uname(1) command (try uname -a)
- Not that the size of the fields in the utsname struct may not be large enough to id a host on a network

To get just a hostname that will identify you on a TCP/IP network, use the Berkeley-dervied:

<pre>#include <unistd.h></unistd.h></pre>									
<pre>int gethostname(char *name,</pre>	int	namelen);							
		Returns:	0	if	OK,	-1	on	error	

- Time is kept in UTC
- Time conversions (timezone, daylight savings time) handled "automatically"
- Time and date kept in a single quantity (time_t)



We can break this time_t value into its components with either of the following:

#include <time.h>

localtime(3) takes into account daylight savings time and the *TZ* environment variable.

The mktime(3) function operates in the reverse direction.

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To output human readable results, use:

#include <time.h>

Lastly, there is a printf(3) like function for times:

#include <time.h>

Homework

Reading:

- Stevens, Chapter 4 and 6
- Falsehoods Programmers believe about time: http://is.gd/yFSYR0

Think about code. You should now be able to implement:

- In(1), mv(1), rm(1), rmdir(1), stat(1)
- date(1), id(1), touch(1), uname(1)

Other:

work on your midterm project!